

Colouring a graph containing no odd circles

Let $G = (N, E)$ be such a graph.

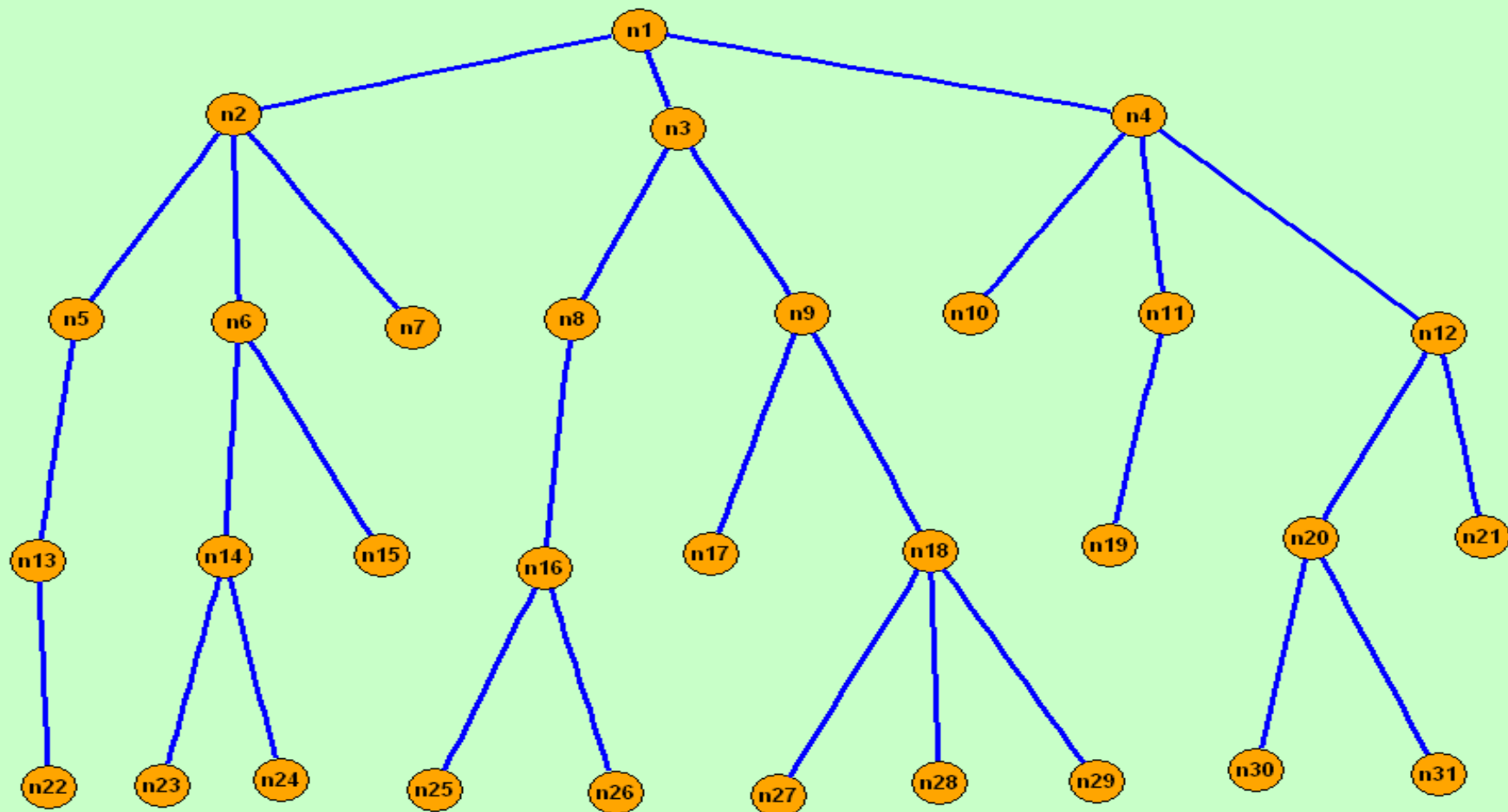
Start from any node v_1 . Assign v_1 level 0.

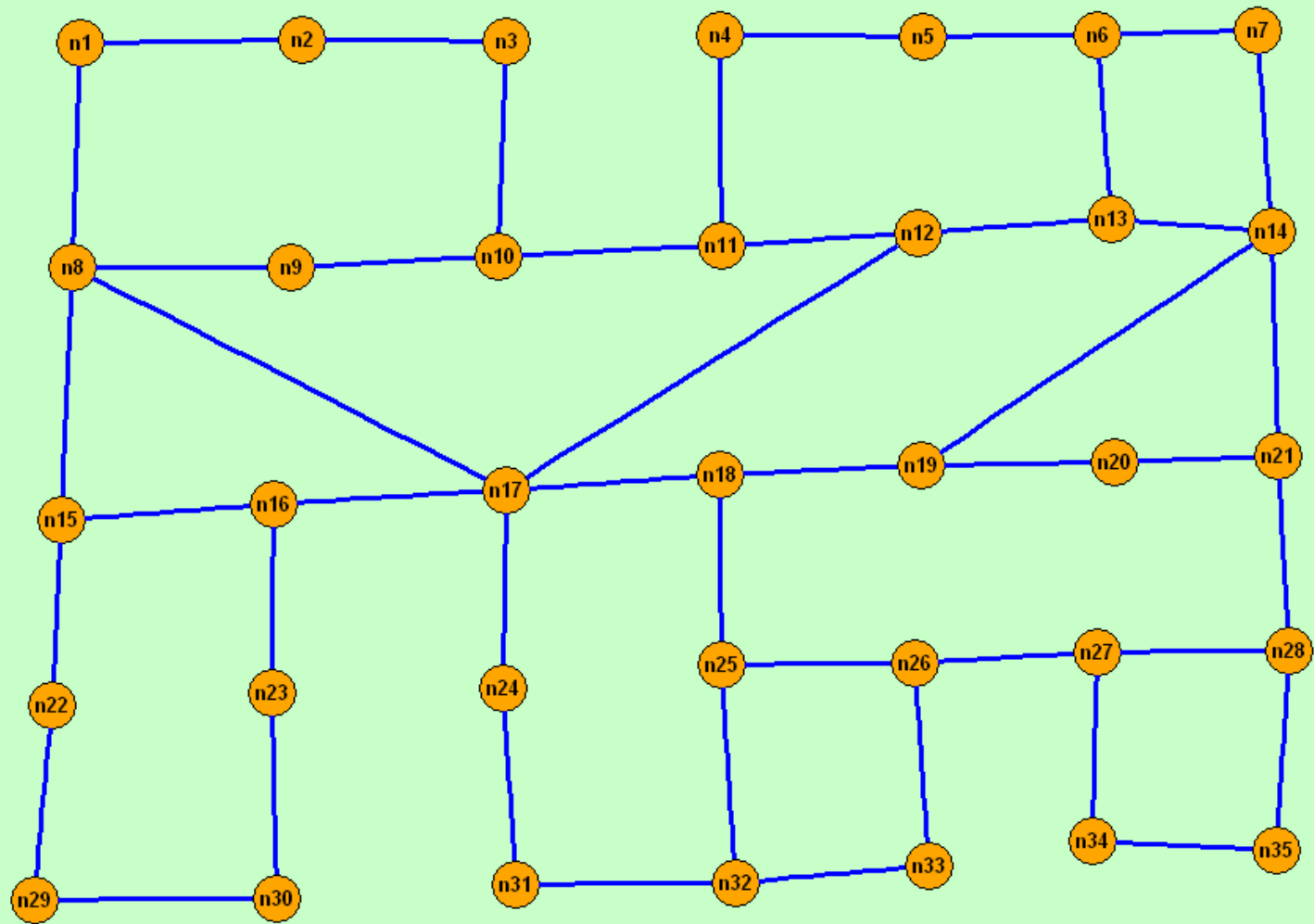
Assign to all nodes adjacent to v_1 level 1.

Partition the set of nodes of G into sets of nodes at different levels.

Let $v_{i_1}, v_{i_2}, \dots, v_{i_r}$ be the nodes at level i . Consider all nodes adjacent to v_{i_1} that are not assigned levels 0, 1, 2, ..., i . Assign these level $i + 1$.

Continue this process until all nodes are assigned a level. Assign now one colour to nodes at odd levels and another colour to all those at an even level.





Approximation Algorithm

No simple characterization is known for graphs that are k -colourable if $k > 2$. In general, it is a hard problem to determine the chromatic number of an arbitrary graph. No algorithm is known that finds a colouring pattern using the fewest possible colours for every input graph. However, there are colouring algorithms that can be used to approximate the best colouring in the sense that they may sometimes use more colours than necessary. The following is an algorithm known as the *largest-first algorithm*:

- 1.order the nodes by their degrees creating a sequence: $\mathbf{v} = v_1, v_2, \dots$ and consider a sequence $\mathbf{c} = c_1, c_2, \dots$ of colours
- 2.assign c_1 to v_1 and keep assigning c_1 to nodes in \mathbf{v} non adjacent to v_1 and previously coloured nodes until you can
- 3.assign c_2 to the node v_i with the greatest degree in \mathbf{v} that has not been assigned a colour and then continue assigning c_2 to others nodes in \mathbf{v} non-adjacent to v_i until you can
- 4.continue this process until all the nodes have been assigned colours

